

Security Aspects of Fuzzy Hashing

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 - ▶ IT-Security
 - ▶ Fuzzy Hashing
- ▶ Since March 2011 Research Student at CASED
 - ▶ Center for Advanced Security Research Darmstadt
- ▶ Publications:
 - ▶ User Survey on Phone Security and Usage (BioSIG Sept. 2010)
 - ▶ Security Aspects of Piecewise Hashing in Computer Forensics (Accepted at IMF Mai 2011)

Motivation [1/2]

- ▶ Main question: Is it possible to identify similar files based on a fingerprint, which depends **only** on the files' byte structure?
- ▶ Cryptographic hash functions follow the avalanche effect: Changing a bit in the input affects $\approx 50\%$ of the output bits
→ no match
- ▶ Fuzzy hashing promises to overcome this problem and discover similarities based on fingerprints.
- ▶ Question addressed in this talk: How reliable are the results of Kornblum's approach for fuzzy hashing with respect to an **active adversary**?

Motivation [2/2] - Applications

1. Forensics (on the file level): Detect similar files
 - ▶ Blacklisting:
 - ▶ Detect manipulated suspicious files
 - ▶ Find fragments of suspicious data
 - ▶ Whitelisting: Find variants of unsuspicious files
2. Biometrics: Template protection
3. Malware: Detect obfuscated malware (e.g. metamorphic malware)
4. Junk mail detection

Agenda

Kornblum's Fuzzy Hashing

Security Aspects

Conclusion

Contact, Discussion

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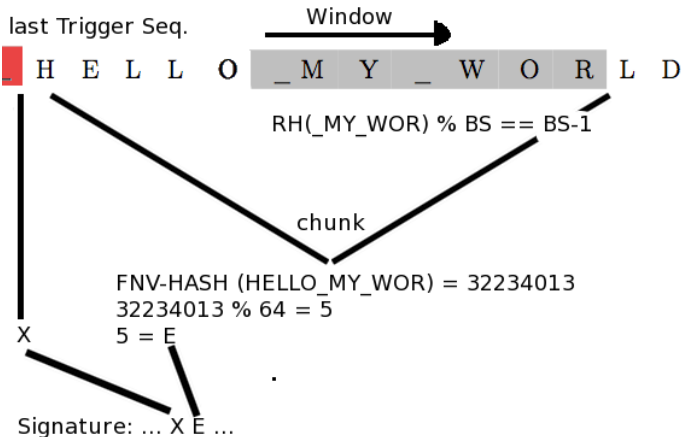
Fuzzy Hashing by Kornblum

- ▶ *Context Triggered Piecewise Hashing* (CTPH)
(software named ssdeep)
- ▶ Developed in 2006 based on spamsum-algorithm from A. Tridgell
- ▶ Key elements:
 - ▶ Block size
 - ▶ Rolling hash
 - ▶ Traditional hash / piecewise hashing
 - ▶ Signature
- ▶ Pioneer: *dcfld*
 - ▶ Blocks had a fixed size
 - ▶ Non-propagation = yes alignment robustness = no

Key Elements

- ▶ Block size: b
 - ▶ $b_{init} = b_{min} \cdot 2^{\lfloor \log_2(\frac{n}{s \cdot b_{min}}) \rfloor}$
- ▶ Rolling Hash at position p in the file:
 - ▶ $r_p = F(n_{p-s+1}, n_{p-s+2}, \dots, n_p)$
 - ▶ Allows to compute r_{p+1} cheaply from r_p by removing the influence of n_{p-s+1} and including the new byte n_{p+1}
- ▶ Traditional Hash / Piecewise Hashing:
 - ▶ Currently, ssdeep makes use of Fowler/Noll/Vo (FNV)
 - ▶ Alternative hash functions are possible (e.g. SHA-1, MD5)

Workflow



Kornblum Signature

- ▶ 2 Signatures:
 - ▶ Signature 1: Using block size b (at most 64 characters)
 - ▶ Signature 2: Using block size $2b$ (at most 32 characters)

- ▶ Sample Kornblum signature of test-file1:

```
1 24:T0tUHZbAzIaFG91Y6pYaK3YKqbaCo/6Pqy45kwUnmJrrevqw+oWluBY5b32TpC0:  
   T0tU5s7ai6ptg7ZNcqMwUArKvqfZlMC0,"/test-file1"
```

Kornblum's Fuzzy Hashing

Security Aspects

Conclusion

Contact, Discussion

Characteristics of Kornblum's Implementation

- ▶ Signature comparison:
 - ▶ Only signatures with the same block size or within a factor of 2 can be compared
 - ▶ A successful match needs at least *one common substring in the signature of length 7*
 - ▶ A signature has at most 64 characters
- ▶ If block size is known, we can calculate trigger sequences:
 - ▶ Easy observation: A trigger sequence for b is also a trigger sequence for all block sizes $\frac{b}{2^k}$
 - ▶ Concatenation of trigger sequences yields signature characters (e.g. _MY_WOR in previous example)
- ▶ Attack type depends on the file syntax:
 - ▶ BMP / ASCII-files can be changed 'everywhere' (easily)
 - ▶ JPG / PDF-files allow a change of header information

Attacks?

What do we like to achieve?

1. **False negatives for blacklisting** → anti-blacklisting
 - ▶ Modified incriminated files are not detected by the blacklist although perceptual identical to the original known-to-be-bad file
2. **False positives for whitelisting** → anti-whitelisting
 - ▶ Incriminated files are modified to get a signature of a known-to-be-good file
 - ▶ Modified incriminated file is perceptual identical to the original known-to-be-bad file

Attacks for Anti-Blacklisting

- ▶ **Blow up a file:** Block size gets larger
- ▶ **Edit trigger seq.:** Block size gets different (unpractical)
- ▶ **Edit between trigger seq.:** Change one byte in every 7th chunk
- ▶ **Adding trigger seq.:** Add several trigger seq. in the beginning of a file
- ▶ No semantic attacks like rotations, colour changes, ...

Anti-Blacklisting: Blow up a File

```
1 $ ls -la hacker_siedlung.jpg
2 -rw-r--r-- 1 user user 68650 2011-02-23 13:57 hacker_siedlung.jpg
3
4 $ dd if=/dev/urandom of=hacker_siedlung.hacked.jpg bs=1 count=280000
5 280000+0 records in
6 280000+0 records out
7 280000 bytes (280 kB) copied, 1.39661 s, 200 kB/s
8
9 $ dd if=hacker_siedlung.jpg of=hacker_siedlung.hacked.jpg conv=notrunc
10 69653+0 records in
11 69653+0 records out
12 69653 bytes (70 kB) copied, 0.20225 s, 344 kB/s
13
14 $ ssdeep -l hacker_siedlung.jpg hacker_siedlung.hacked.jpg
15 ssdeep,1.0--blocksize:hash:hash,filename
16 1536:FLVoUaX+ns+6iAuLNdElzt/CclGbn20CFN8DXg1BSXHaL++:
17 F3l6ew33lG20MBSXa6+,"hacker_siedlung.jpg"
18 6144:F6j0MBEjZML1AecfyqefFgQ5wDg+b7LQ7vZ0ubiPZ:
19 F40Mq6i8qefFgU1Tsub6Z,"hacker_siedlung.hacked.jpg"
```

Attacks for Anti-Whitelisting

- ▶ **Edit between trigger seq.:** Change one byte in each chunk
- ▶ **Adding trigger seq.:** Add several trigger seq. in the beginning
- ▶ Difference: Adding information vs. editing information
- ▶ More computational power than for anti-blacklisting

Example: Editing Between Trigger Sequences

H E L L O _ M Y _ W O R L D

Example: Adding Trigger Sequences

- ▶ File need to be changed in the beginning
- ▶ One may use **global trigger sequences**:

Trigger Sequence	Base64 Char.	Trigger Sequence	Base64 Char.
AAAD?Hp	9	AAAV?Hf	l
AAAD?Og	v	AAAf?Ft	p
AAAD?QI	7	AAAr?xj	V
AAAJ?MW	P	AAAx?Fj	1
AAAJ?PJ	F	AAAx?OC	n
AAAJ?VO	Z	AAAx?tx	5

Table 3.1.: Sample pre-computed global trigger sequences and their corresponding Base64 signature characters

- ▶ Example: Insertion of concatenation of trigger sequences
AAAD?HpAAAD?OgAAAD?QIAAAJ?MWAAAJ?PJAAAJ?VO yields
Kornblum's signature: 9v7PFZ

Kornblum's Fuzzy Hashing

Security Aspects

Conclusion

Contact, Discussion

Summary

- ▶ CTPH from Kornblum does not withstand an active adversary with respect to
 - ▶ blacklisting
 - ▶ whitelisting
- ▶ Doubtful if piecewise hashing can fulfill the expectations of fuzzy hashing
 - ▶ Typically it is possible to flip one bit in each chunk
- ▶ In order to create a viable new fuzzy hash function, it will be necessary to find different approaches

Future Work

- ▶ Conduct a study if CTPH is applicable in forensics
- ▶ Clear definition of:
 - ▶ What we expect from a fuzzy hash function?
 - ▶ What is a metric for similarity?
- ▶ Find a more general approach, which also addresses images, videos, ... not only txt files
- ▶ Proof if this might be possible on byte level
- ▶ Otherwise new techniques might be needed:
 - ▶ E.g. first extract features then hash (e.g. FFT for images)

Thank you for your attention!

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