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Sebastian Altmeyer, Claire Burguière, Jan Reineke

Angaben zur Veröffentlichung / Publication details:

Altmeyer, Sebastian, Claire Burguière, and Jan Reineke. 2009. "Cache related preemption delay for set-associative caches: resilience analysis." In *AVACS Workshop, Oldenburg, Germany, 2009*. Augsburg: Universität Augsburg. https://www.rw.cdl.uni-saarland.de/people/reineke/private/talks/resilienceAnalysisAVACS09.pdf.

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Cache Related Preemption Delay for Set-Associative Caches

Resilience Analysis

Sebastian Altmeyer, Claire Burguière, Jan Reineke

AVACS Workshop, Oldenburg 2009

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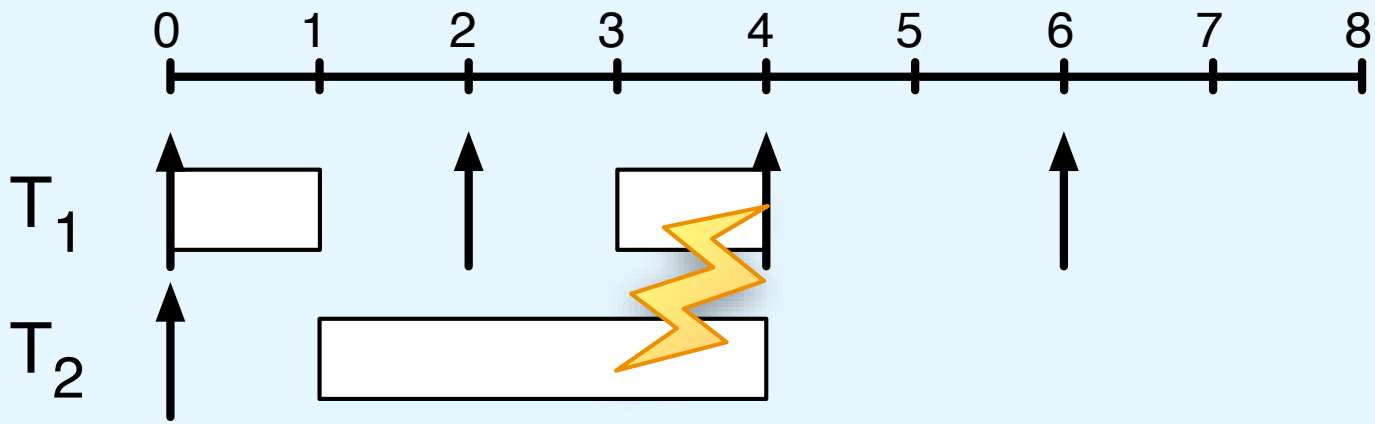
Why use preemptive scheduling?

- Preemption often increases schedulability of task sets.
- Tasks with short deadlines are often unschedulable non-preemptively.

Example

Given: Two periodic tasks T_1 and T_2, with periods $P_1 = 2, P_2 = 8$, deadlines $D_1 = P_1, D_2 = P_2$, and execution times $C_1 = 1, C_2 = 3$.



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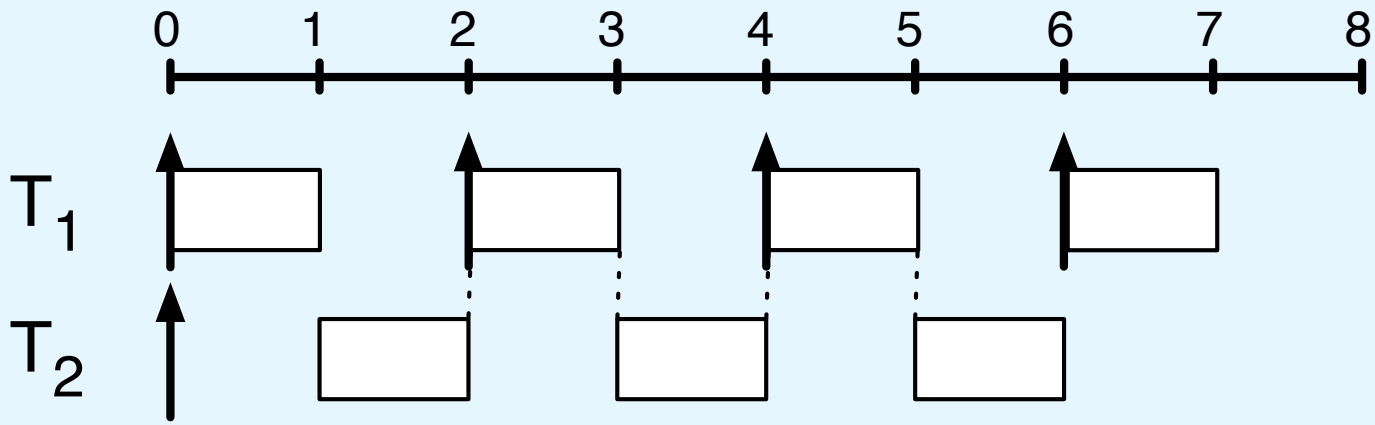
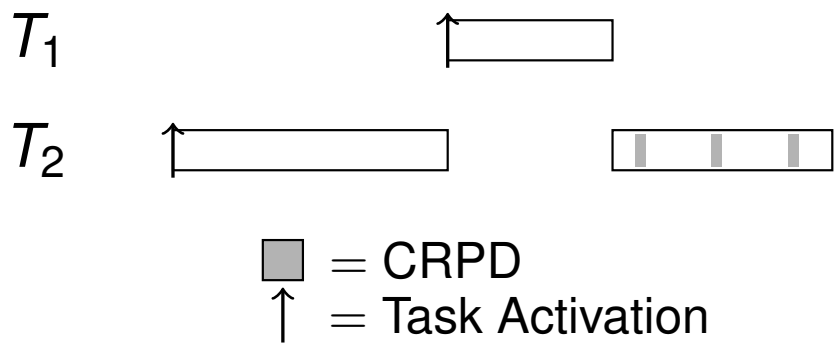


Preemption does not come for free!

- The preempting task “disturbs” the state of performance-enhancing features like caches and pipelines.
- Once the preempted task resumes its execution, the disturbance may cause additional *cache misses*.
- The additional execution time due to additional cache misses is known as the *cache-related preemption delay* (CRPD).



How to take preemption cost into account?

Where to account for preemption cost?

- Integrate into WCET Analysis: [**?**]
  - assume cache misses everywhere
  - very pessimistic but easy to use in schedulability analysis
- WCET Analysis + CRPD Analysis: [**?**]
  - $WCET_{bound} + n \cdot CRPD_{bound} \geq$ execution time of task with up to n preemptions
  - more precise but not supported by many schedulability analyses

CRPD for set-associative caches - LRU

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- CRPD computation:
  - Preempted task: Useful Cache Blocks (UCB)
  - Preempting task: Evicting Cache Blocks (ECB)
- CRPD from UCB **and** ECB:
  - Previous combination rather imprecise

  $\Rightarrow$ Some UCBs remain useful under preemption

Useful Cache Block - [?]

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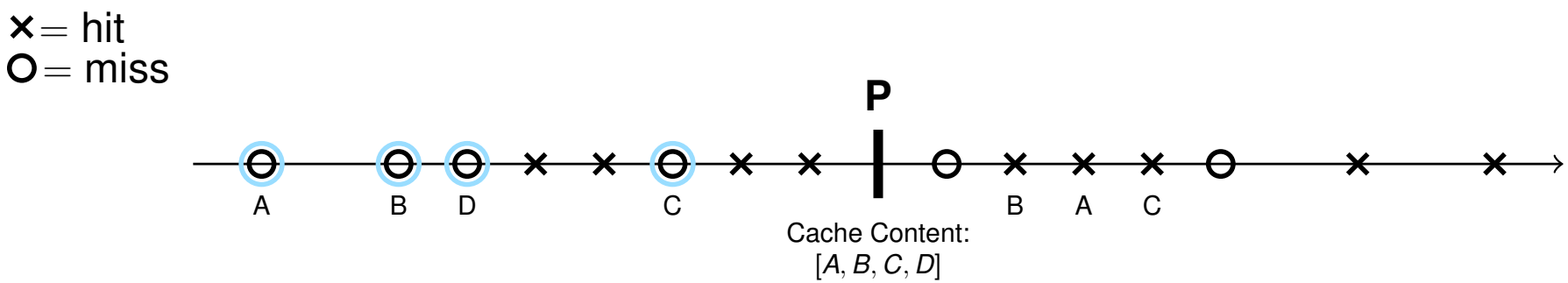
Definition (Useful Cache Block)

A memory block m at program point P is called a useful cache block, if

a) *m may be cached at P*

b) *m may be reused at program point P′ that may be reached from P with no eviction of m on this path.*



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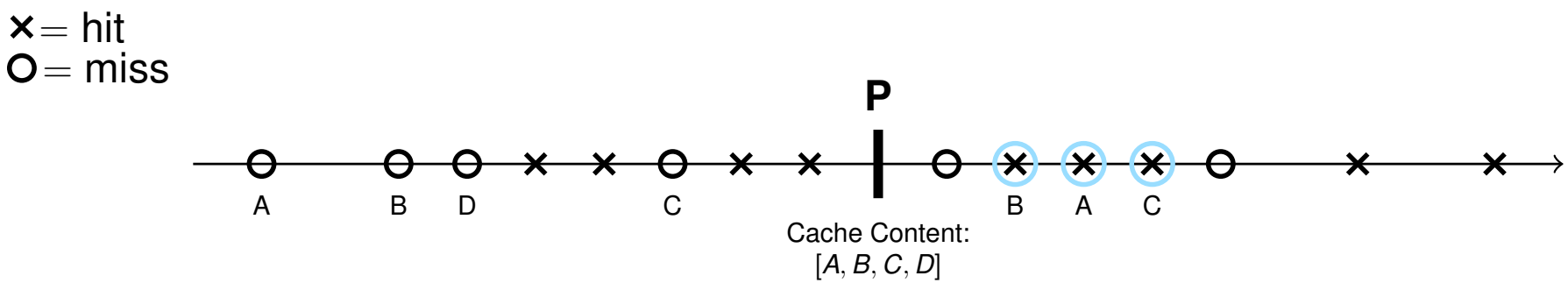

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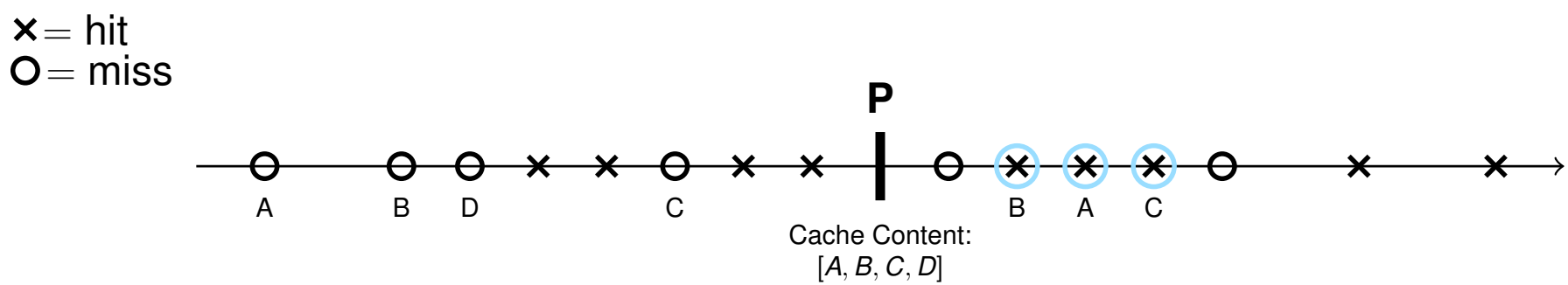

$$\mathrm{CRPD}_{\mathrm{UCB}} = \sum_{s=1}^{c} \mathrm{CRPD}^{s}_{\mathrm{UCB}}$$

$$\mathrm{CRPD}^{s}_{\mathrm{UCB}} = \mathrm{BRT} \times \mathit{min}(|\mathrm{UCB}(s)|, n)$$

n = associativity
BRT = Block Reload Time

Evicting Cache Blocks [?]

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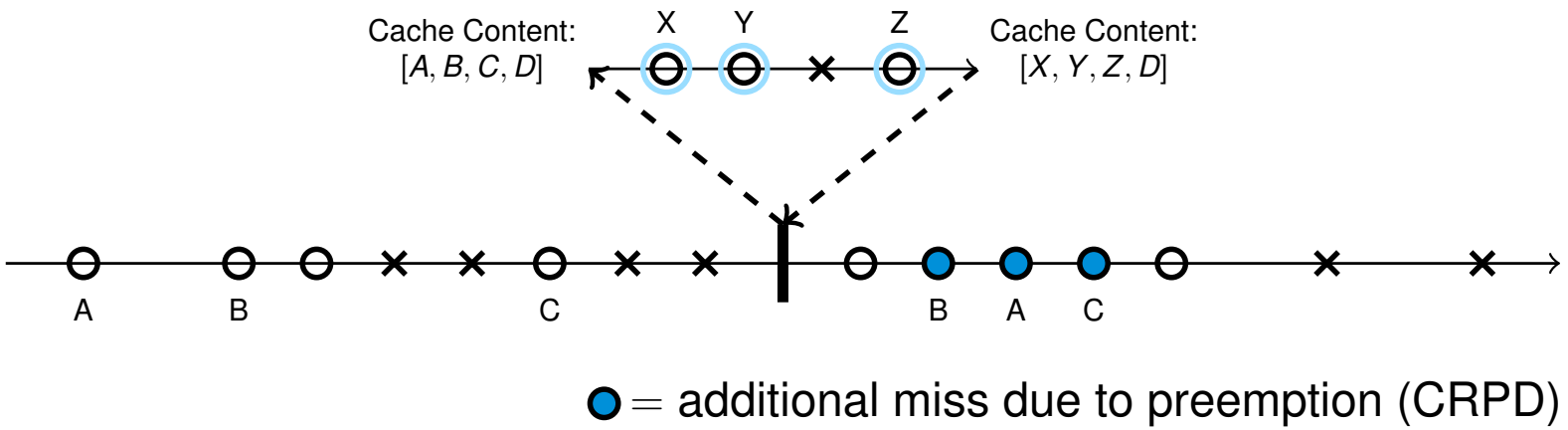
Definition (Evicting Cache Blocks (ECB))

A memory block of the preempting task is called an evicting cache block, if it may be accessed during the execution of the preempting task.



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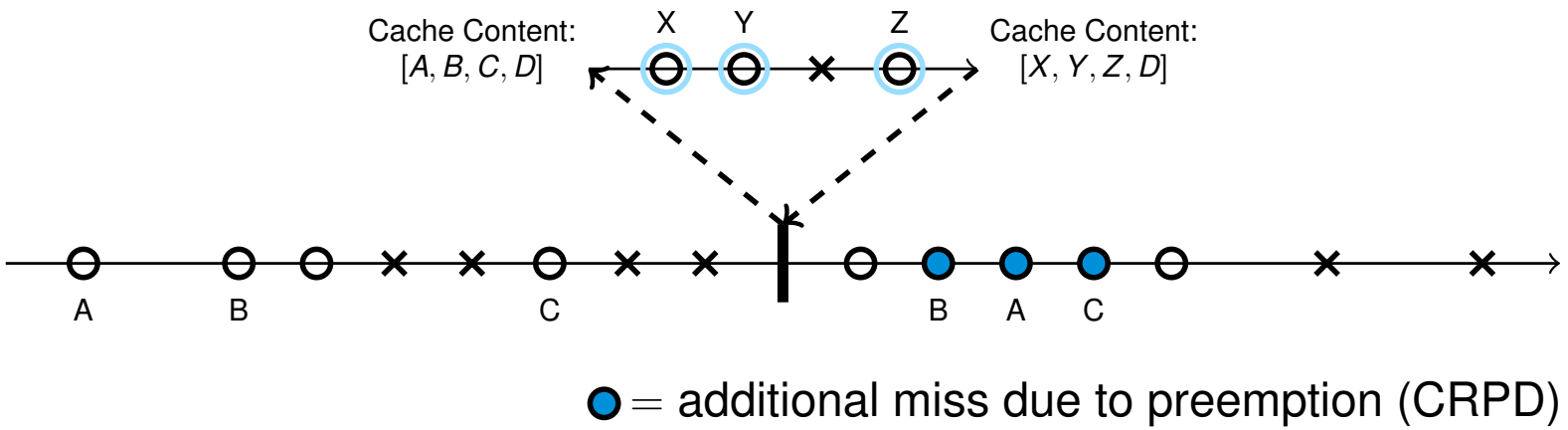

$$\mathrm{CRPD}^{s}_{\mathrm{ECB}} = \begin{cases} 0 & \text{if } \mathrm{ECB}(s) = \emptyset \\ \mathrm{BRT} \times n & \text{otherwise} \end{cases}$$

Impact of the preempting task on the preempted task

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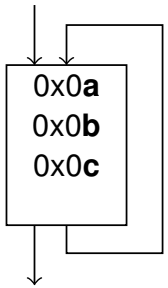
CRPD (using UCB and ECB)

$$CRPD_{UCB\&ECB} = \sum_{s=1}^{c} min(CRPD^{s}_{UCB}, CRPD^{s}_{ECB})$$

Impact of the preempting task on the preempted task: Example



$$[c, b, a, x] \xrightarrow{a} [a, c, b, x] \xrightarrow{b} [b, a, c, x] \xrightarrow{c} [c, b, a, x]$$ no misses

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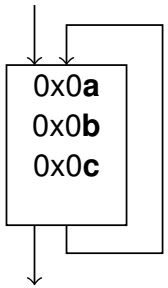

ECBs $= \{\mathbf{e}\}$

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- $\text{CRPD}_{\text{UCB}} \Rightarrow |\text{UCB}| = 3$
- $\text{CRPD}_{\text{ECB}} \Rightarrow n = 4$
- $\text{CRPD}_{\text{UCB\&ECB}} = min(\text{CRPD}_{\text{UCB}}, \text{CRPD}_{\text{ECB}}) \Rightarrow 3$
  - Overestimation: number of additional misses $= 0 < 3$

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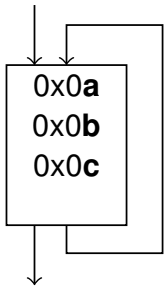

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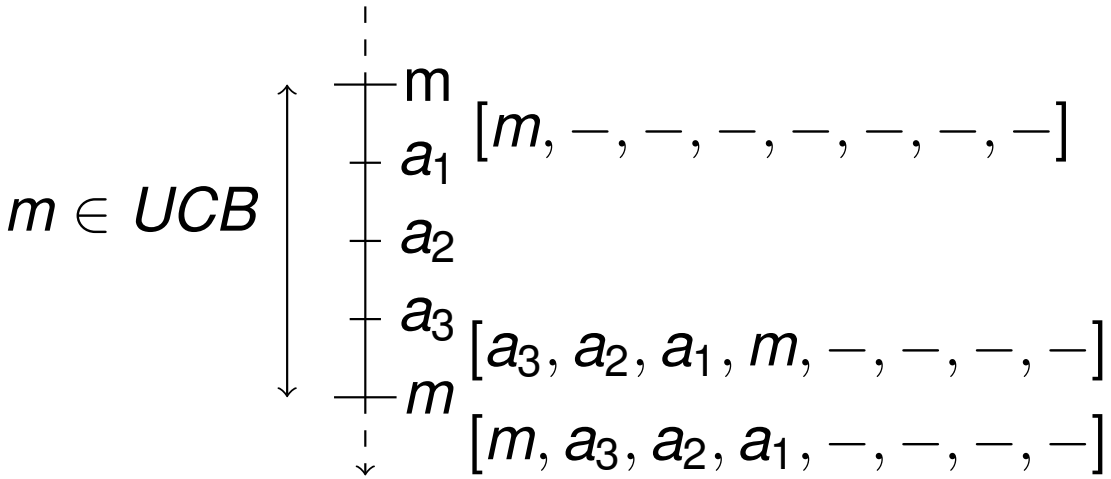
  - Overestimation: number of additional misses $= 0 < 3$
- Why?
  - $|\text{ECB}|$ to evict a UCB $= 2$, but
  - $|\text{ECB}| = 1$
  - A single ECB is not sufficient to evict a UCB.

Combining UCB and ECB: Refinement

Determining $max|\mathrm{ECB}|$, such that no additional cache miss occur



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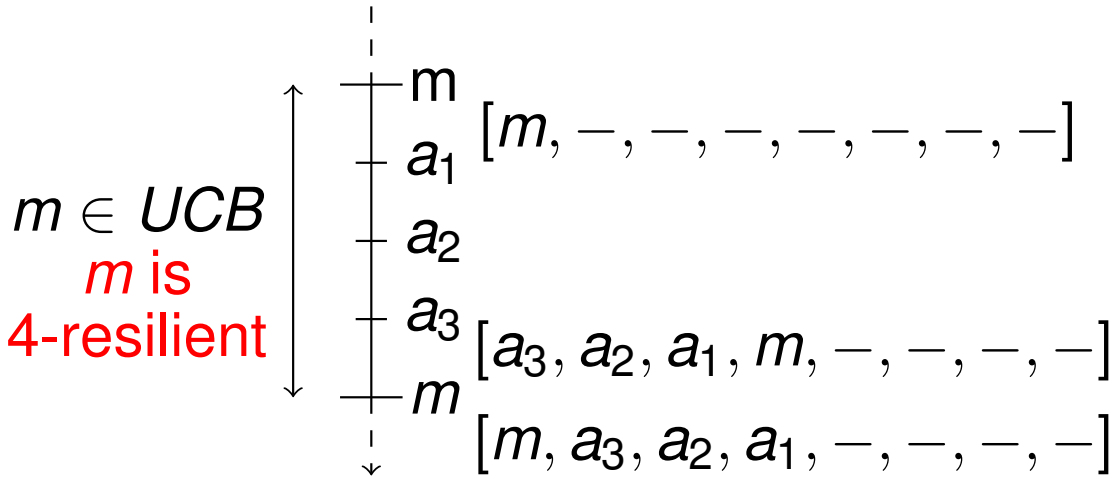

Resilience analysis

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Definition (l-Resilience)

A memory block m is called l-resilient at program point P,
if all possible next accesses to m

- *that would be hits without preemption,*
- *would still be hits in case of a preemption at P with l accesses.*

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- No UCB is *n*-resilient, i.e., no UCB remains useful after a preemption with *n* ECBs.
- Each $(l+1)$-resilient UCB is also l-resilient.
- Each UCB is at least 0-resilient.

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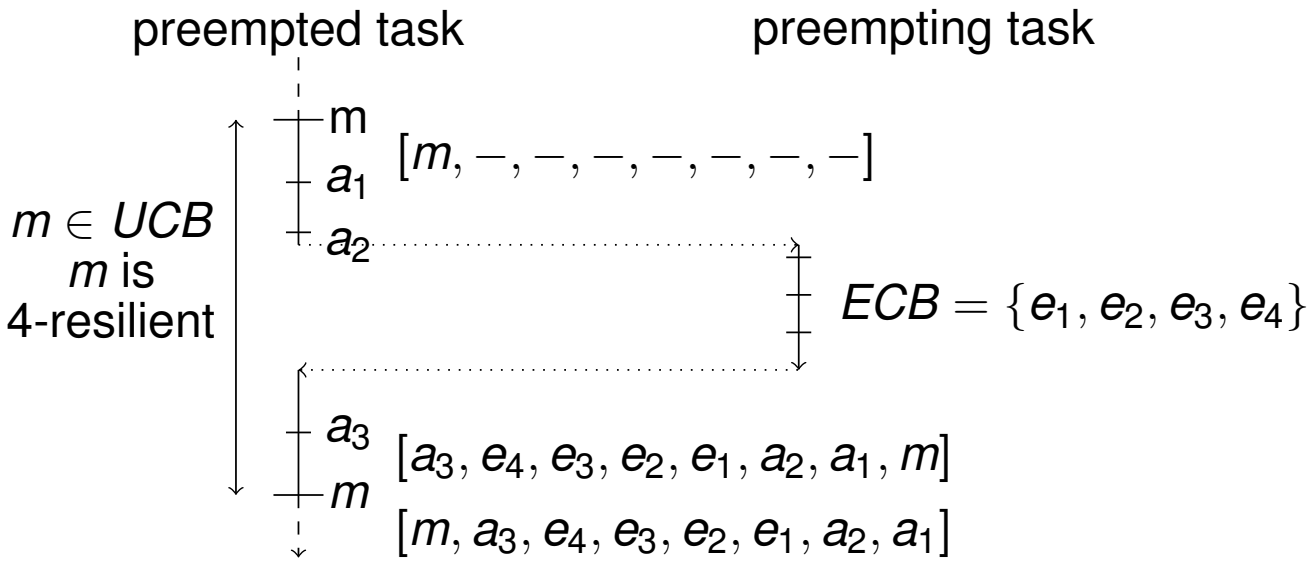

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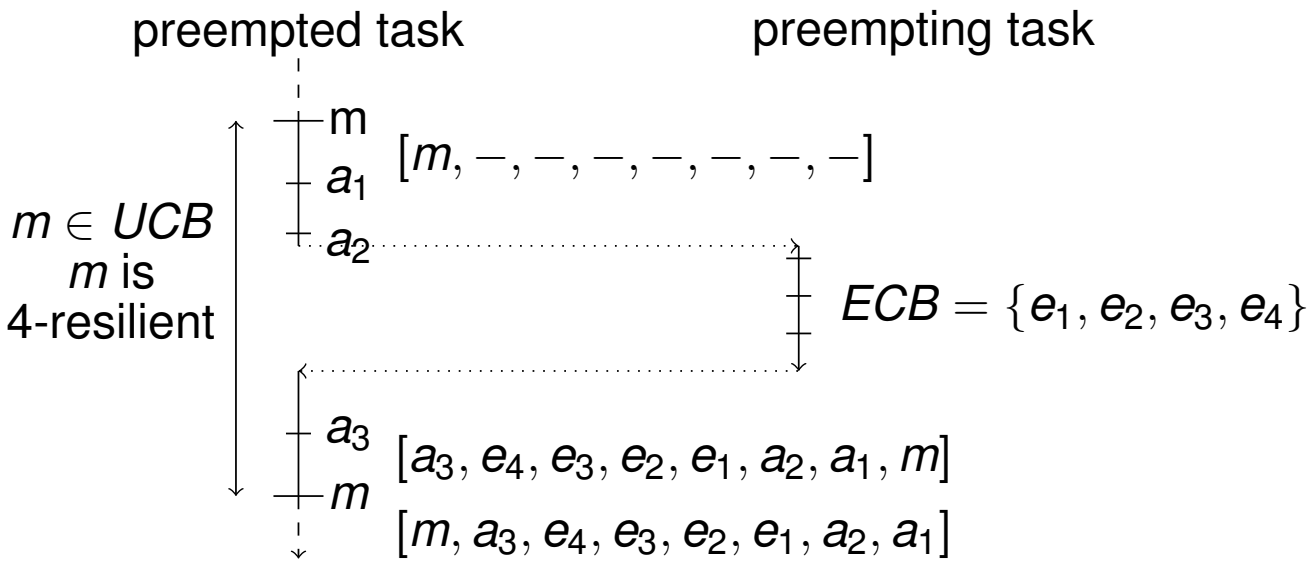

In general: if $|\text{ECB}| \leq l$ then the UCB is not evicted

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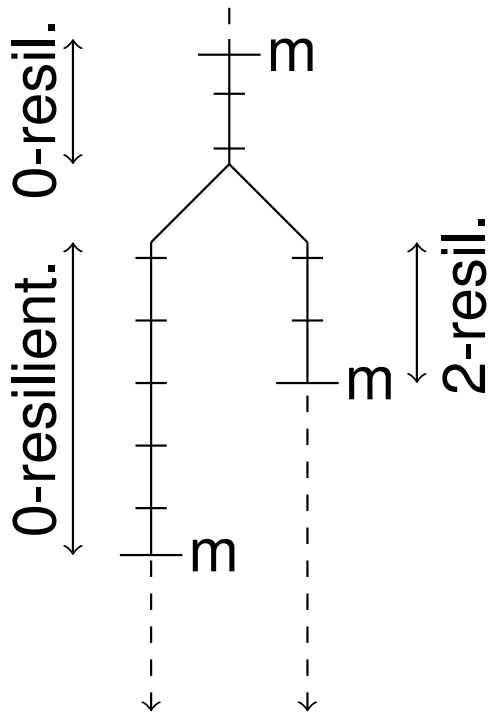

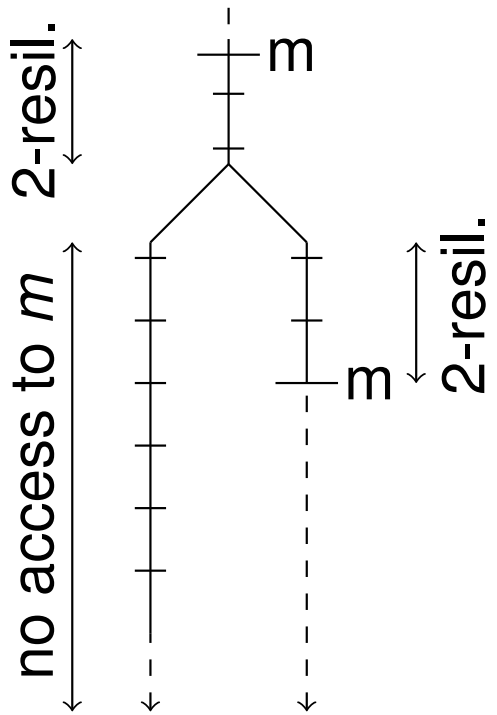

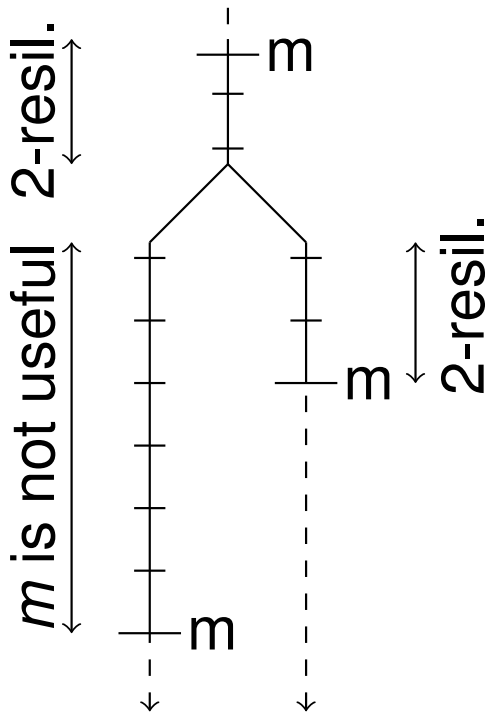

Bounding the CRPD using Resilience

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CRPD (Combining UCB and ECB by using Resilience)

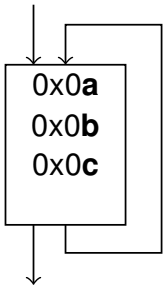
$$\text{blocks contributing to CRPD}: \quad \Big| \underbrace{UCB}_{\text{useful}} \setminus \underbrace{\{m \mid m \text{ is } |ECB|\text{-resilient}\}}_{\text{remain useful}} \Big|$$

Bounding the CRPD using Resilience

CRPD (Combining UCB and ECB by using Resilience)

$$CRPD \leq BRT \times |\overbrace{\underbrace{UCB}_{\text{useful}} \setminus \{m \mid m \text{ is } |ECB|\text{-resilient}\}}^{\text{blocks contributing to CRPD}}|$$

Bounding the CRPD using Resilience: Example



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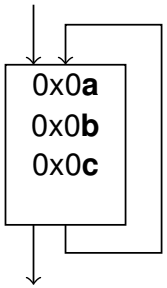

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- $|\text{ECB}| = 1$
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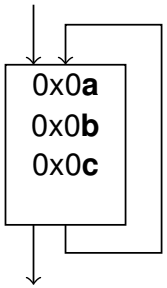

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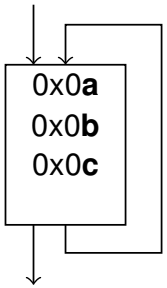

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Conclusions

- Preemptive scheduling:
  - sometimes necessary
  - but not for free: CRPD
- UCB and ECB analyses:
  - pessimistic overapproximation of the CRPD
- Resilience analysis:
  - determining the set of UCBs that remain useful under preemption
  - increase precision
  - implemented as two simple data-flow analyses:
    - ★ similar to UCB analysis for LRU
    - ★ currently in the phase of evaluation

Further reading

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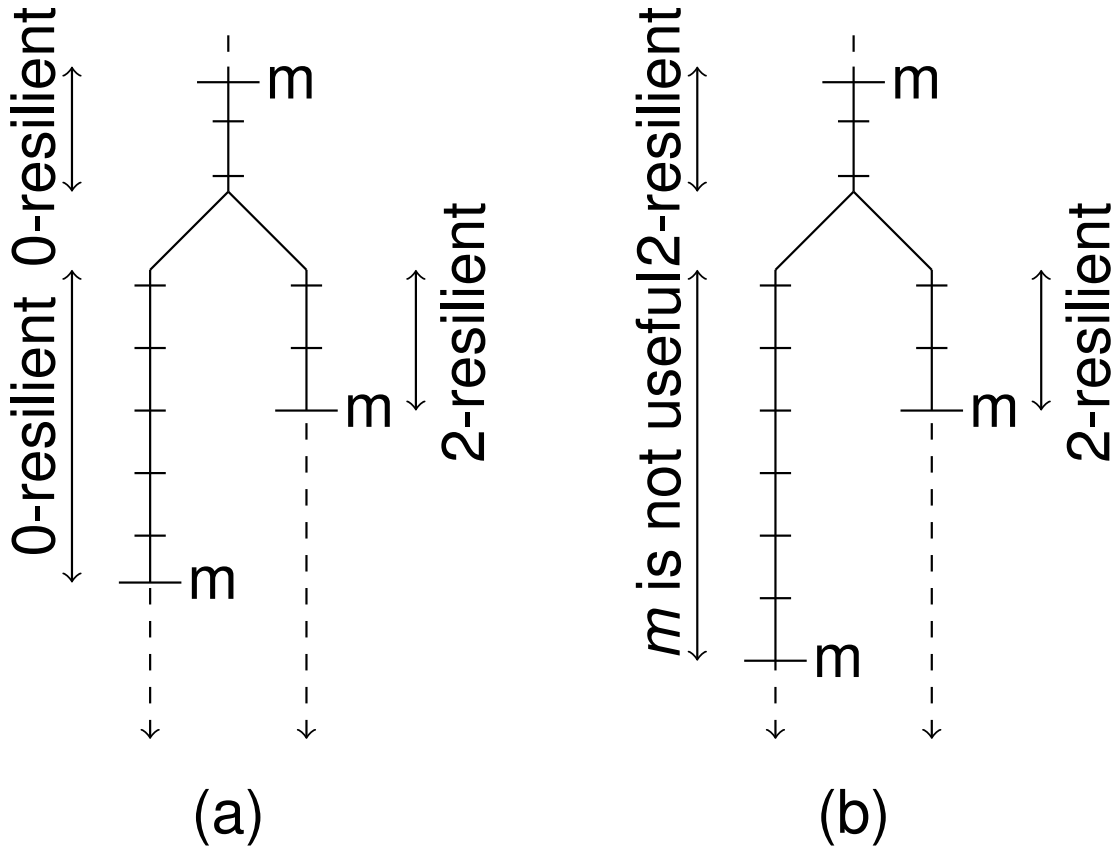
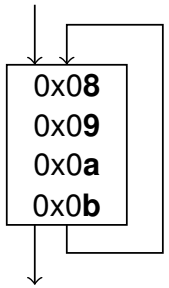
/-resilience analysis



CPRD using ECB: Pitfall



ECBs $= \{\mathbf{e}\}$

$[b, a, 9, 8] \xrightarrow{8} [8, b, a, 9] \xrightarrow{9} [9, 8, b, a] \xrightarrow{a} [a, 9, 8, b] \xrightarrow{b} [b, a, 9, 8]$ 0 misses

$[\mathbf{e}, b, a, 9] \xrightarrow{8^*} [8, \mathbf{e}, b, a] \xrightarrow{9^*} [9, 8, \mathbf{e}, b] \xrightarrow{a^*} [a, 9, 8, \mathbf{e}] \xrightarrow{b^*} [b, a, 9, 8]$ 4 misses

- $|\mathrm{UCB}(s)| = 4$
- $|\mathrm{ECB}(s)| = 1$
- $n = 4$
- number of additional misses$= 4$

Upper-bound on the CRPD - direct-mapped caches

- using UCB [**?**]:

$$\mathrm{CRPD}_{\mathrm{UCB}} = \mathrm{BRT} \cdot |\{s_i \mid \exists m \in \mathrm{UCB} : m \bmod c = s_i\}|$$

- using ECB [**?**]:

$$\mathrm{CRPD}_{\mathrm{ECB}} = \mathrm{BRT} \cdot |\{s_i \mid \exists m \in \mathrm{ECB} : m \bmod c = s_i\}|$$

- using UCB and ECB [**?**, **?**]:

$$\begin{aligned}\mathrm{CRPD}_{\mathrm{UCB\&ECB}} = \mathrm{BRT} \cdot |\{s_i \mid\ & \exists m \in \mathrm{UCB} : m \bmod c = s_i \\ & \wedge \exists m' \in \mathrm{ECB} : m' \bmod c = s_i\}|\end{aligned}$$

CRPD for FIFO: Pitfalls

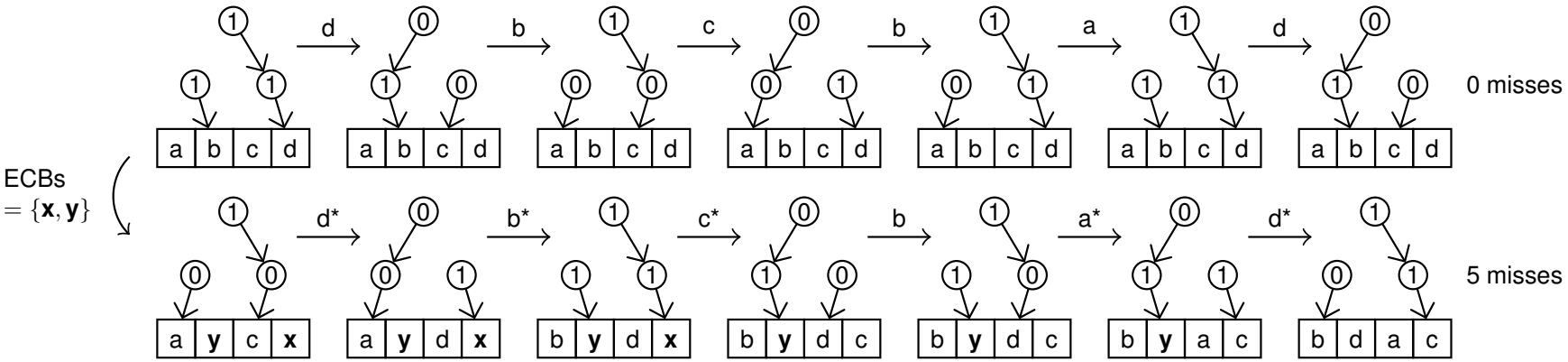
$$\text{ECBs} = \{\mathbf{x}\} \begin{cases} [b,a] \xrightarrow{a} [b,a] \xrightarrow{e^*} [e,b] \xrightarrow{b} [e,b] \xrightarrow{c^*} [c,e] \xrightarrow{e} [c,e] & \text{2 misses} \\ [\mathbf{x},b] \xrightarrow{a^*} [a,\mathbf{x}] \xrightarrow{e^*} [e,a] \xrightarrow{b^*} [b,e] \xrightarrow{c^*} [c,b] \xrightarrow{e^*} [e,c] & \text{5 misses} \end{cases}$$

CRPD for FIFO: Pitfalls

$$\text{ECBs} = \{\mathbf{x}\} \begin{cases} [b,a] \xrightarrow{a} [b,a] \xrightarrow{e^*} [e,b] \xrightarrow{b} [e,b] \xrightarrow{c^*} [c,e] \xrightarrow{e} [c,e] & \text{2 misses} \\ [\mathbf{x},b] \xrightarrow{a^*} [a,\mathbf{x}] \xrightarrow{e^*} [e,a] \xrightarrow{b^*} [b,e] \xrightarrow{c^*} [c,b] \xrightarrow{e^*} [e,c] & \text{5 misses} \end{cases}$$

- $|\mathrm{UCB}(s)| = 2$
- $|\mathrm{ECB}(s)| = 1$
- $n = 2$
- But: number of additional misses$= 3$

CRPD for PLRU: Pitfalls



- $|\mathrm{UCB}(s)| = 4$
- $|\mathrm{ECB}(s)| = 2$
- $n = 4$
- But: number of additional misses$= 5$